

Generalized Measures of Edge Fault Tolerance in (n, k) -star Graphs*

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Abstract

This paper considers a kind of generalized measure $\lambda_s^{(h)}$ of fault tolerance in the (n, k) -star graph $S_{n,k}$ for $2 \leq k \leq n-1$ and $0 \leq h \leq n-k$, and determines $\lambda_s^{(h)}(S_{n,k}) = \min\{(n-h-1)(h+1), (n-k+1)(k-1)\}$, which implies that at least $\min\{(n-k+1)(k-1), (n-h-1)(h+1)\}$ edges of $S_{n,k}$ have to remove to get a disconnected graph that contains no vertices of degree less than h . This result shows that the (n, k) -star graph is robust when it is used to model the topological structure of a large-scale parallel processing system.

Keywords: Combinatorics, fault-tolerant analysis, (n, k) -star graphs, edge-connectivity, h -super edge-connectivity

1 Introduction

It is well known that interconnection networks play an important role in parallel computing/communication systems. An interconnection network can be modeled by a graph $G = (V, E)$, where V is the set of processors and E is the set of communication links in the network. The connectivity $\lambda(G)$ of a graph G is an important measurement for fault-tolerance of the network, and the larger $\lambda(G)$ is, the more reliable the network is.

A subset of vertices B of a connected graph G is called a *edge-cut* if $G - B$ is disconnected. The *edge connectivity* $\lambda(G)$ of G is defined as the minimum cardinality over all edge-cuts of G . Because λ has many shortcomings, one proposes the concept of the h -super edge connectivity of G , which can measure fault tolerance of an interconnection network more accurately than the classical connectivity λ .

A subset of vertices B of a connected graph G is called an *h -super edge-cut*, or *h -edge-cut* for short, if $G - B$ is disconnected and has the minimum degree at least h . The *h -super edge-connectivity* of G , denoted by $\lambda_s^{(h)}(G)$, is defined as the minimum cardinality over all h -edge-cuts of G . It is clear that, if $\lambda_s^{(h)}(G)$ exists, then

$$\lambda(G) = \lambda_s^{(0)}(G) \leq \lambda_s^{(1)}(G) \leq \lambda_s^{(2)}(G) \leq \cdots \leq \lambda_s^{(h-1)}(G) \leq \lambda_s^{(h)}(G).$$

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For any graph G and integer h , determining $\lambda_s^{(h)}(G)$ is quite difficult. In fact, the existence of $\lambda_s^{(h)}(G)$ is an open problem so far when $h \geq 1$. Some results have been obtained on $\lambda_s^{(h)}$ for particular classes of graphs and small h 's (see Section 16.7 in [5]).

This paper is concerned about $\lambda_s^{(h)}$ for the (n, k) -star graph $S_{n,k}$. In h -super connectivity, several authors have done some work. For $k = n - 1$, $S_{n,n-1}$ is isomorphic to a star graph S_n . Akers and Krishnamurthy [1] determined $\lambda(S_n) = n - 1$ for $n \geq 2$ and $\lambda_s^{(1)}(S_n) = 2n - 4$ for $n \geq 3$. In this paper, we show the following result.

Theorem 1.1 *If $2 \leq k \leq n - 1$ and $0 \leq h \leq n - k$, then*

$$\lambda_s^{(h)}(S_{n,k}) = \begin{cases} (n - h - 1)(h + 1) & \text{for } h \leq k - 2 \text{ and } h \leq \frac{n}{2} - 1, \\ (n - k + 1)(k - 1) & \text{otherwise.} \end{cases}$$

This result implies that at least $\min\{(n - k + 1)(k - 1), (n - h - 1)(h + 1)\}$ edges of $S_{n,k}$ have to remove to get a disconnected graph that contains no vertices of degree less than h . The proof of this result is in Section 3. In Section 2, we recall the structure of $S_{n,k}$ and some lemmas used in our proofs.

2 Definitions and lemmas

For given integer n and k with $1 \leq k \leq n - 1$, let $I_n = \{1, 2, \dots, n\}$ and $P(n, k) = \{p_1 p_2 \dots p_k : p_i \in I_n, p_i \neq p_j, 1 \leq i \neq j \leq k\}$, the set of k -permutations on I_n . Clearly, $|P(n, k)| = n! / (n - k)!$.

Definition 2.1 *The (n, k) -star graph $S_{n,k}$ is a graph with vertex-set $P(n, k)$. The adjacency is defined as follows: a vertex $p = p_1 p_2 \dots p_i \dots p_k$ is adjacent to a vertex*

- (a) $p_i p_2 \dots p_{i-1} p_1 p_{i+1} \dots p_k$, where $2 \leq i \leq k$ (swap p_1 with p_i).
- (b) $\alpha p_2 p_3 \dots p_k$, where $\alpha \in I_n \setminus \{p_i : 1 \leq i \leq k\}$ (replace p_1 by α).

The vertices of type (a) are referred to as *swap-neighbors* of p and the edges between them are referred to as *swap-edge* or *i-edges*. The vertices of type (b) are referred to as *unswap-neighbors* of p and the edges between them are referred to as *unswap-edges*. Clearly, every vertex in $S_{n,k}$ has $k - 1$ swap-neighbors and $n - k$ unswap-neighbors. Usually, if $x = p_1 p_2 \dots p_k$ is a vertex in $S_{n,k}$, we call p_i the *i-th bit* for each $i \in I_k$.

The (n, k) -star graph $S_{n,k}$ is proposed by Chiang and Chen [4]. Some nice properties of $S_{n,k}$ are compiled by Cheng and Lipman (see Theorem 1 in [2]).

Lemma 2.2 *$S_{n,k}$ is $(n - 1)$ -regular $(n - 1)$ -connected.*

Lemma 2.3 *For any $\alpha = p_1 p_2 \dots p_{k-1} \in P(n, k - 1)$ ($k \geq 2$), let $V_\alpha = \{p\alpha : p \in I_n \setminus \{p_i : i \in I_{k-1}\}\}$. Then the subgraph of $S_{n,k}$ induced by V_α is a complete graph of order $n - k + 1$, denoted by K_{n-k+1}^α .*

Let $S_{n-1,k-1}^{t:i}$ denote a subgraph of $S_{n,k}$ induced by vertices with the t -th bit i for $2 \leq t \leq k$. The following lemma is a slight modification of the result of Chiang and Chen [4].

Lemma 2.4 For a fixed integer t with $2 \leq t \leq k$, $S_{n,k}$ can be decomposed into n subgraphs $S_{n-1,k-1}^{t:i}$, which is isomorphic to $S_{n-1,k-1}$, for each $i \in I_n$. Moreover, there are $\frac{(n-2)!}{(n-k)!}$ independent swap-edges between $S_{n-1,k-1}^{t:i}$ and $S_{n-1,k-1}^{t:j}$ for any $i, j \in I_n$ with $i \neq j$.

Since $S_{n,1} \cong K_n$, we only consider the case of $k \geq 2$ in the following discussion.

Lemma 2.5 If $2 \leq k \leq n-1$ and $0 \leq h \leq n-k$, then

$$\lambda_s^{(h)}(S_{n,k}) \leq \begin{cases} (n-h-1)(h+1) & \text{for } h \leq \frac{n}{2} - 1, \\ (n-k+1)(k-1) & \text{otherwise.} \end{cases}$$

Proof. By our hypothesis of $h \leq n-k$, for any $\alpha \in P(n, k-1)$, we can choose a subset $X \subseteq V(K_{n-k+1}^\alpha)$ such that $|X| = h+1$. Then the subgraph of K_{n-k+1}^α induced by X is a complete graph K_{h+1} . Let B be the set of incident edges with and not within X . Since $S_{n,k}$ is $(n-1)$ -regular and K_{h+1} is h -regular, we have that

$$|B| = (n-h-1)(h+1).$$

Clearly, B is an edge-cut of $S_{n,k}$. Let x be any vertex in $S_{n,k} - X$, and $d(x)$ denote the number of edges incident with x in $S_{n,k} - X$. In order to prove that B is an h -edge-cut, we only need to show $d(x) \geq h$. Note that X is contained in $S_{n-1,k-1}^i$ and edges between $S_{n-1,k-1}^i$ and $S_{n-1,k-1}^j$ are independent for any $i, j \in I_n$ with $i \neq j$ by Lemma 2.4. If x is in $S_{n-1,k-1}^i - K_{n-k+1}$ or is in $S_{n-1,k-1}^j$ with $i \neq j$, then $d(x) \geq n-2 \geq n-k \geq h$. For $x \in V(K_{n-k+1} - X)$, if exists, then $d(x) = n-1-|X| = n-h-2 \geq h$ for $h \leq \frac{n}{2} - 1$. Therefore, B is an h -edge-cut of $S_{n,k}$, and so

$$\lambda_s^{(h)}(S_{n,k}) \leq |B| = (n-h-1)(h+1) \text{ for } h \leq \frac{n}{2} - 1.$$

If $h \geq \frac{n}{2}$, we choose $X = V(K_{n-k+1}^\alpha)$. Then $|B| = (n-k+1)(k-1)$. For any x in $S_{n-1,k-1}^i - X$ or $S_{n-1,k-1}^j$ with $i \neq j$, we have $d(x) \geq n-2 \geq n-k \geq h$. Thus, B is an h -edge-cut of $S_{n,k}$, and so

$$\lambda_s^{(h)}(S_{n,k}) \leq |B| = (n-k+1)(k-1) \text{ for } h \geq \frac{n}{2}.$$

The lemma follows. ■

Corollary 2.6 $\lambda_s^{(h)}(S_{n,2}) = n-1$ for $0 \leq h \leq n-2$.

Proof. On the one hand, $\lambda_s^{(h)}(S_{n,2}) \leq n-1$ by Lemma 2.5 when $k=2$. On the other hand, $\lambda_s^{(h)}(S_{n,2}) \geq \lambda(S_{n,2}) = n-1$ by Lemma 2.2. ■

The following lemma shows the relations between $(n-h-1)(h+1)$ and $(n-k+1)(k-1)$.

Lemma 2.7 For $2 \leq k \leq n-1, 0 \leq h \leq n-k$, let

$$\psi(h, k) = \min\{(n-h-1)(h+1), (n-k+1)(k-1)\}. \quad (2.1)$$

If $h \leq \frac{n}{2} - 1$, then

$$\psi(h, k) = \begin{cases} (n-h-1)(h+1) & \text{if } 0 \leq h \leq k-2; \\ (n-k+1)(k-1) & \text{if } h \geq k-1. \end{cases}$$

Proof. Let $f(x) = (n - x)x$, then $\psi(h, k) = \min\{f(h + 1), f(k - 1)\}$. It can be easily checked that $f(x)$ is a convex function on the interval $[0, n]$, the maximum value is reached at $x = \frac{n}{2}$. Thus, $f(x)$ is an increasing function on the interval $[0, \frac{n}{2}]$.

If $0 \leq h \leq k - 2$, then $h + 1 \leq k - 1$. Since $h \leq n - k$, $h + 1 \leq n - k + 1$. Clearly, $\max\{k - 1, n - k + 1\} \leq \frac{n}{2}$. Thus, when $h \leq \frac{n}{2} - 1$, $f(h + 1) \leq f(k - 1) = f(n - k + 1)$, and so $\psi(h, k) = f(h + 1) = (n - h - 1)(h + 1)$.

If $h \geq k - 1$, then $k - 1 < h + 1 \leq \frac{n}{2}$. Thus, $f(k - 1) < f(h + 1)$, and so $\psi(n, k) = f(k - 1) = (n - k + 1)(k - 1)$.

The lemma follows. \blacksquare

To state and prove our main results, we need some notations. Let B be a minimum h -edge-cut of $S_{n,k}$. Clearly, $S_{n,k} - B$ has exactly two connected components. Let X and Y be two vertex-set of two connected components of $S_{n,k} - B$. For a fixed $t \in I_k \setminus \{1\}$ and any $i \in I_n$, let

$$\begin{aligned} X_i &= X \cap V(S_{n-1,k-1}^{t:i}), \\ Y_i &= Y \cap V(S_{n-1,k-1}^{t:i}), \\ B_i &= B \cap E(S_{n-1,k-1}^{t:i}) \text{ and} \\ B_{ij} &= B \cap E(S_{n-1,k-1}^{t:i}, S_{n-1,k-1}^{t:j}), \end{aligned} \quad (2.2)$$

and let

$$\begin{aligned} J &= \{i \in I_n : X_i \neq \emptyset\}, \\ J' &= \{i \in J : Y_i \neq \emptyset\} \text{ and} \\ T &= \{i \in I_n : Y_i \neq \emptyset\}. \end{aligned} \quad (2.3)$$

Lemma 2.8 *Let B be a minimum h -edge-cut of $S_{n,k}$ and X be the vertex-set of a connected component of $S_{n,k} - B$. If $3 \leq k \leq n - 1$ and $1 \leq h \leq n - k$ then, for any $t \in I_k \setminus \{1\}$,*

- (a) B_i is an $(h - 1)$ -edge-cut of $S_{n-1,k-1}^{t:i}$ for any $i \in J'$,
- (b) $\lambda_s^{(h)}(S_{n,k}) \geq |J'| \lambda_s^{(h-1)}(S_{n-1,k-1})$,

Proof. (a) By the definition of J' , B_i is an edge-cut of $S_{n-1,k-1}^{t:i}$ for any $i \in J'$. For any vertex x in $S_{n-1,k-1}^{t:i} - B_i$, since x has degree at least h in $S_{n,k} - S$ and has exactly one neighbor outsider $S_{n-1,k-1}^{t:i}$, x has degree at least $h - 1$ in $S_{n,k}^{t:i} - B_i$. This fact shows that B_i is an $(h - 1)$ -edge-cut of $S_{n-1,k-1}^{t:i}$ for any $i \in J'$.

(b) By the assertion (a), we have $|B_i| \geq \lambda_s^{(h-1)}(S_{n-1,k-1})$, and so

$$\lambda_s^{(h)}(S_{n,k}) = |B| \geq \sum_{i \in J'} |B_i| \geq |J'| \lambda_s^{(h-1)}(S_{n-1,k-1}).$$

the lemma follows. \blacksquare

3 Proof of Theorem 1.1

Proof. By Lemma 2.5 and Lemma 2.7, we only need to prove that, for $2 \leq k \leq n - 1$ and $0 \leq h \leq n - k$,

$$\lambda_s^{(h)}(S_{n,k}) \geq \begin{cases} (n - h - 1)(h + 1) & \text{for } h \leq k - 2 \text{ and } h \leq \frac{n}{2} - 1, \\ (n - k + 1)(k - 1) & \text{otherwise.} \end{cases} \quad (3.1)$$

Let $\omega(h, k) = \max\{(n - h - 1)(h + 1), (n - k + 1)(k - 1)\}$.

We proceed by induction on $k \geq 2$ and $h \geq 0$. The inequality (3.1) is true for $k = 2$ and any h with $0 \leq h \leq n - 2$ by Corollary 2.6. The inequality (3.1) is also true for $h = 0$ and any k with $2 \leq k \leq n - 1$ since $\lambda_s^{(0)}(S_{n,k}) = \lambda(S_{n,k}) = n - 1$. Assume the induction hypothesis for $k - 1$ with $k \geq 3$ and for $h - 1$ with $h \geq 1$, that is,

$$\lambda_s^{(h-1)}(S_{n-1,k-1}) \geq \begin{cases} (n-h)h & \text{for } h \leq k-3 \text{ and } h \leq \frac{n-1}{2}, \\ (n-k+2)(k-2) & \text{otherwise.} \end{cases} \quad (3.2)$$

Let B be a minimum h -edge-cut of $S_{n,k}$ and X be the vertex-set of a minimum connected component of $S_{n,k} - B$. By Lemma 2.5, we have

$$|B| \leq \omega(h, k). \quad (3.3)$$

Use notations defined in (2.2) and (2.3). Choose $t \in I_k \setminus \{1\}$ such that $|J|$ is as large as possible. For each $i \in I_n$, we write $S_{n-1,k-1}^i$ for $S_{n-1,k-1}^{t:i}$ for short.

We first show $|J| = 1$. Suppose to the contrary $|J| \geq 2$. We will deduce contradictions by considering three cases depending on $|J'| = 0$, $|J'| = 1$ or $|J'| \geq 2$.

Case 1. $|J'| = 0$,

In this case, $X_i \neq \emptyset$ and $Y_i = \emptyset$ for each $i \in J$, that is, $J \cap T = \emptyset$. By $|J| \geq 2$ and the minimality of X , $|T| \geq 2$. Assume $\{i_1, i_2\} \subseteq J$ and $\{i_3, i_4\} \in T$. By Lemma 2.4, there are $\frac{(n-2)!}{(n-k)!}$ independent swap-edges between $S_{n-1,k-1}^{i_1}$ (resp. $S_{n-1,k-1}^{i_2}$) and $S_{n-1,k-1}^{i_3}$ (resp. $S_{n-1,k-1}^{i_4}$), all of which are contained in B . Since $J \cap T = \emptyset$, we have that

$$|B| \geq 4 \frac{(n-2)!}{(n-k)!}.$$

For $k = 3$,

$$|B| \geq 4 \frac{(n-2)!}{(n-k)!} \geq 4(n-2) > 2(n-2)$$

Combining Lemma 2.5 with Lemma 2.7 yields $|B| \leq \lambda_s^{(h)}(S_{n,3}) \leq 2(n-2)$, a contradiction. For $k \geq 4$, it is easy to check that

$$\begin{aligned} |B| &\geq 4 \frac{(n-2)!}{(n-k)!} \geq 4(n-2)(n-3) = (2n-4)(2n-6) \\ &> \max\{(n-h-1)(h+1), (n-k+1)(k-1)\} \\ &= \omega(h, k), \end{aligned}$$

which contradicts the inequality (3.3).

Case 2. $|J'| = 1$,

Without loss of generality, assume $J' = \{1\}$. By Lemma 2.8 (a), B_1 is an $(h-1)$ -edge-cut of $S_{n-1,k-1}^1$.

By $|J| \geq 2$, there exists an $i \in J - J'$ such that $X_i = V(S_{n-1,k-1}^i)$. By the minimality of X , there exists some $j \in T - J'$ such that $Y_j = V(S_{n-1,k-1}^j)$. By Lemma 2.4, there are $\frac{(n-2)!}{(n-k)!}$ independent swap-edges between $S_{n-1,k-1}^i$ and $S_{n-1,k-1}^j$, thus $|B_{ij}| = \frac{(n-2)!}{(n-k)!} \geq n-2$. By (3.2), we consider the following two cases.

If $\lambda_s^{(h-1)}(S_{n-1,k-1}) \geq (n-h)h$, then

$$\begin{aligned} |B| &\geq |B_1| + |B_{ij}| \\ &\geq (n-h)h + (n-2) \\ &> (n-h-1)h + (n-h-1) \\ &= (n-h-1)(h+1), \end{aligned}$$

If $\lambda_s^{(h-1)}(S_{n-1,k-1}) \geq (n-k+2)(k-2)$, then

$$\begin{aligned} |B| &\geq |B_1| + |B_{ij}| \\ &\geq (n-k+2)(k-2) + (n-2) \\ &> (n-k+2)(k-2) + (n-k+2) \\ &> (n-k+1)(k-1). \end{aligned}$$

Therefore, we have $|B| > \omega(h, k)$, which contradicts the inequality (3.3).

Case 3. $|J'| \geq 2$.

By Lemma 2.8 (b) and (3.2), we consider the following two cases.

If $\lambda_s^{(h-1)}(S_{n-1,k-1}) \geq (n-h)h$, then

$$\begin{aligned} |B| &\geq |J'| \lambda_s^{(h-1)}(S_{n-1,k-1}) \\ &\geq 2(n-h)h \geq (n-h)h + (n-h) \\ &> (n-h-1)(h+1), \end{aligned}$$

If $\lambda_s^{(h-1)}(S_{n-1,k-1}) \geq (n-k+2)(k-2)$, then

$$\begin{aligned} |B| &\geq |J'| \lambda_s^{(h-1)}(S_{n-1,k-1}) \\ &\geq 2(n-k+2)(k-2) \\ &\geq (n-k+2)(k-2) + (n-k+2) \\ &> (n-k+1)(k-1). \end{aligned}$$

Therefore, we have $|B| > \omega(h, k)$, which contradicts the inequality (3.3).

Thus, we have $|J| = 1$. By the choice of t , the i -th bits of all vertices in X are same for each $i = 2, 3, \dots, k$, and so X is a complete graph. Thus, we have that

$$\lambda_s^{(h)}(S_{n,k}) = |B| = (n - |X|)|X|$$

Since $h+1 \leq |X| \leq n-k+1$ and $f(x) = (n-x)x$ is a convex function on the interval $[0, n]$, we have that

$$\lambda_s^{(h)}(S_{n,k}) = |B| = (n - |X|)|X| \geq \psi(h, k),$$

where $\psi(h, k)$ is defined in (2.1).

If $h \leq \frac{n}{2} - 1$, using Lemma 2.7, we have

$$\lambda_s^{(h)}(S_{n,k}) \geq \psi(h, k) = \begin{cases} (n-h-1)(h+1) & \text{if } 0 \leq h \leq k-2; \\ (n-k+1)(k-1) & \text{if } h \geq k-1. \end{cases} \quad (3.4)$$

If $h \geq \frac{n}{2}$, we have $X = V(K_{n-k+1})$. Otherwise, there exists some $x \in V(K_{n-k+1} - X)$ such that

$$h \leq d(x) = n-1-|X| \leq n-h-2,$$

which implies $h \leq \frac{n}{2} - 1$, a contradiction. Therefore, we have $|X| = n-k+1$, and

$$\lambda_s^{(h)}(S_{n,k}) = |B| = (n - |X|)|X| = (n-k+1)(k-1) \text{ for } h \geq \frac{n}{2}. \quad (3.5)$$

Combining (3.4) with (3.5) yields (3.1). By the induction principle, the theorem follows. \blacksquare

As we have known, when $k = n-1$, $S_{n,n-1}$ is isomorphic to the star graph S_n . Akers and Krishnamurthy [1] determined $\lambda(S_n)$ and $\lambda_s^{(1)}(S_n)$, which can be obtained from our result by setting $k = n-1$ and $h = 0, 1$, respectively.

Corollary 3.1 (Akers and Krishnamurthy [1]) $\lambda(S_n) = n-1$ for $n \geq 2$ and $\lambda_s^{(1)}(S_n) = 2n-4$ for $n \geq 3$.

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